# ICFP M2 - Statistical Physics 2 - TD n ${ }^{\circ} 5$ Random XORSAT problems 

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We shall consider in this problem random systems of linear equations, also known as XORSAT problems, denoted $F$. We recall their definition given during the lectures :

- the degrees of freedom are $N$ Boolean variables, $\underline{x}=\left(x_{1}, \ldots, x_{N}\right) \in\{0,1\}^{N}$
- they have to obey $M$ linear constraints of the form

$$
\begin{equation*}
x_{i_{a}^{1}}+x_{i_{a}^{2}}+\cdots+x_{i_{a}^{k}}=y_{a} \bmod 2, \tag{1}
\end{equation*}
$$

where $a=1, \ldots, M$ indexes the various equations, $k \geq 3$ is an integer defining the number of variables involved in each equation, $\left\langle i_{a}^{1}, \ldots, i_{a}^{k}\right\rangle$ is a $k$-uplet of distinct indices in $\{1, \ldots, N\}$, and $y_{a} \in\{0,1\}$ fixes the right hand side of the equation.

- such a formula is said to be satisfiable if there is a configuration $\underline{x}$ that verifies all the equations simultaneously, unsatisfiable otherwise.
- a random ensemble of formulas is defined very easily by generating the $M$ equations independently, choosing for each of them a $k$-uplet $\left\langle i_{a}^{1}, \ldots, i_{a}^{k}\right\rangle$ uniformly at random among the $\binom{N}{k}$ possible ones, and $y_{a}=0$ or 1 with probability $1 / 2$.
Using the Gaussian elimination algorithm one can determine whether a given formula is satisfiable or not in polynomial time. Repeating this process a large number of times one can easily obtain a numerical estimate of the probability $P_{\text {sat }}(\alpha, N)$ that a random formula $F$ with $N$ variables and $M=\alpha N$ equations is satisfiable :


These curves, obtained for $k=3$, suggest that a phase transition occurs in the thermodynamic limit ( $N, M \rightarrow$ $\infty$ with $\alpha=M / N$ fixed) for $\alpha$ around 0.92 . Indeed, there exists a threshold $\alpha_{\text {sat }}$ (that depends on $k$ ) such that

$$
\lim _{N \rightarrow \infty} P_{\mathrm{sat}}(\alpha, N)=\left\{\begin{array}{ll}
1 & \text { if } \alpha<\alpha_{\mathrm{sat}}  \tag{2}\\
0 & \text { if } \alpha>\alpha_{\mathrm{sat}}
\end{array} .\right.
$$

## 1 Bounds on $\alpha_{\text {sat }}$

We recall a result obtained in TD2 : for a random variable $Z$ that takes non-negative integer values,

$$
\begin{equation*}
\frac{\mathbb{E}[Z]^{2}}{\mathbb{E}\left[Z^{2}\right]} \leq \mathbb{P}[Z>0] \leq \mathbb{E}[Z] \tag{3}
\end{equation*}
$$

these two inequalities being called the second and first moment method, respectively.
We shall use these inequalities with $Z$ the number of solutions of a random XORSAT formula with $N$ variables and $M$ equations constructed as above.

1. Compute $\mathbb{E}[Z]$, and deduce that $\alpha_{\text {sat }} \leq 1$.
2. Show that

$$
\begin{equation*}
\mathbb{E}\left[Z^{2}\right]=2^{N} \sum_{D=0}^{N}\binom{N}{D}\left(\frac{1}{2} \sum_{\substack{l=0 \\ l \text { even }}}^{k} \frac{\binom{D}{l}\binom{N-D}{k-l}}{\binom{N}{k}}\right)^{M} \tag{4}
\end{equation*}
$$

3. Study the asymptotic behavior of this expression in the thermodynamic limit, and conclude that

$$
\lim _{N \rightarrow \infty} \frac{1}{N} \ln \left(\frac{\mathbb{E}[Z]^{2}}{\mathbb{E}\left[Z^{2}\right]}\right)=\inf _{d \in[0,1]} g(\alpha, d), \quad \text { with } g(\alpha, d)=\ln 2+d \ln d+(1-d) \ln (1-d)-\alpha \ln \left(1+(1-2 d)^{k}\right) .
$$

4. Draw the shape of the function $g$ as a function of $d$ for increasing values of $\alpha$. Conclude that there exists a value $\alpha_{\mathrm{lb}}>0$ (equal to 0.889 for $k=3$ ) such that for $\alpha<\alpha_{\mathrm{lb}}$, the first term in (3) is not exponentially small. A more detailed analysis of (4) shows that in this case it actually goes to 1 . Conclude that $\alpha_{\mathrm{lb}} \leq \alpha_{\mathrm{sat}} \leq 1$.

## 2 Leaf removal procedure

The bounds on $\alpha_{\text {sat }}$ obtained above are not tight (i.e. $\alpha_{\mathrm{lb}}<1$ ) because of the potentially huge fluctuations of $Z$, that can cause its average $\mathbb{E}[Z]$ to be much larger than its typical value. These fluctuations can be reduced by concentrating on a well-chosen subformula, as explained now.

1. Suppose that $F$ contains a leaf, i.e. a variable that appears in a single equation, and denote $F^{\prime}$ the system of equations obtained by removing this equation. Show that $F$ is satisfiable if and only if $F^{\prime}$ is satisfiable.

This leaf removal procedure can be iterated as long as leaves are present. Two cases can occur : either the formula is completely emptied by this procedure, or there remains a non-trivial subset of $F$, called its core, in which every variable appears in at least two equations. We call $N_{\text {core }}$ and $M_{\text {core }}$ the number of variables and equations of the core formula, and display on the curves below the fraction $f_{\text {core }}=N_{\text {core }} / N$ of variables in the core and the density $\alpha_{\text {core }}=M_{\text {core }} / N_{\text {core }}$ of equations it contains.


These curves demonstrate a core percolation transition at $\alpha_{\mathrm{d}}=0.818$ (for $k=3$ ), and show that the density $\alpha_{\text {core }}$ crosses 1 at $\alpha_{*}=0.918$ (for $k=3$ ).
2. A calculation (not required here) shows that

$$
\begin{equation*}
f_{\text {core }}=1-e^{-\alpha k \phi^{k-1}}-\alpha k \phi^{k-1} e^{-\alpha k \phi^{k-1}}, \quad \frac{1}{N} M_{\text {core }}=\alpha \phi^{k} \tag{5}
\end{equation*}
$$

where $\phi=\phi(\alpha, k)$ is the largest solution in $[0,1]$ of the equation

$$
\begin{equation*}
\phi=1-e^{-\alpha k \phi^{k-1}} . \tag{6}
\end{equation*}
$$

Study graphically this equation, show that for $k \geq 3$ the transition at $\alpha_{\mathrm{d}}$ is discontinuous, and study the behavior of $\phi$ for $\alpha \rightarrow \alpha_{\mathrm{d}}^{+}$.
3. Explain why $\alpha_{*}$ is an improved upperbound on $\alpha_{\text {sat }}$. It turns out that the fluctuations of the core are much smaller than that of the full formula, hence actually $\alpha_{\mathrm{sat}}=\alpha_{*}$.

